Positive singular Artin monoids inject into their singular Artin monoids

Noelle Antony

Abstract

We define a map \( \varphi \) from a given singular Artin monoid to its corresponding group algebra and show that \( \varphi \) is faithful precisely when Birman’s conjecture – that the desingularisation map is injective – is true. We discover combinatorial properties of \( \varphi \) which apply to positive singular Artin monoids of any, not necessarily finite, type and infer that \( \varphi \) is injective on pairs of words for which a common multiple exists. Furthermore \( \varphi \) preserves the Intermediate Property, derived by Corran, in the sense that if \( \varphi(\tau_i U) = \varphi(\tau_j V) \) then \( m_{ij} = 2 \) or \( i = j \). From this it follows that \( \varphi \) is injective for some classes of monoids which include singular Artin monoids of type \( I_2(p) \), generalising a result of East. Finally, we invoke Paris’ discovery, that all Artin monoids inject into their groups, as well as our combinatorial discoveries, to deduce that all positive singular Artin monoids embed into their corresponding singular Artin monoids, extending a result of Corran.

AMS Subject Classification: 20F36, 20M05, 20M30.

1 Preliminaries

We begin with a finite indexing set \( I \), and we let \( \Gamma^M \) be a complete labelled graph with \( n \) vertices in one-to-one correspondence with a finite indexing set \( I \) and with edge labels from the set \( \{3, 4, 5, \ldots, \infty\} \). For \( i \neq j \) let \( m_{ij} \) denote the label of the edge between nodes \( i \) and \( j \), or set \( m_{ij} = 2 \) if there is no such edge. Let \( \langle xy \rangle^m \) denote the alternating product \( xyx \ldots \) of length \( m \). Put \( S = \{\sigma_i \mid i \in I\} \), \( T = \{\tau_i \mid i \in I\} \), and let \( S^{-1} = \{\sigma_i^{-1} \mid i \in I\} \), the set of formal inverses of \( S \). If \( X \) is a set then \( X^* \) refers to the free monoid generated by \( X \). The Artin group of type \( M \), denoted \( G_M \), is the group generated by \( S \) subject to the relations

\[
\langle \sigma_i \sigma_j \rangle^{m_{ij}} = \langle \sigma_j \sigma_i \rangle^{m_{ij}} \quad \text{for} \ i, j \in I, \ \text{and} \ \ m_{ij} \neq \infty
\]

denoted \( \mathcal{R}_1 \) and called the braid relations. The positive Artin monoid of type \( M \), \( G_M^+ \), is the monoid generated by \( S \) subject to the braid relations \( \mathcal{R}_1 \). The Coxeter group of type \( M \) is the group generated by \( S \) subject to the preceding relations and the relations \( \sigma_i^2 = 1 \) for every \( i \) in \( I \). In this way we see that Coxeter groups arise as quotient groups of Artin groups. If the Coxeter group of type \( M \) is finite, then \( M \) is said to be of finite type.
Except when explicitly stated, we assume throughout this paper that \( M \) is of any type. The singular Artin monoid of type \( M \), denoted \( SG_M \), is the monoid generated by \( S \cup S^{-1} \cup T \) and has as its defining relations \( \mathcal{R}_1 \), the free group relations \( \sigma_i \sigma_i^{-1} = \sigma_i^{-1} \sigma_i = 1 \) and the relations \( \mathcal{R}_2 \) listed below,

\[
\tau_i \sigma_i = \sigma_i \tau_i \quad \text{for all } i \in I,
\]

\[
\tau_i \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} = \begin{cases} 
\langle \sigma_j \sigma_i \rangle^{m_{ij}-1} \tau_j & \text{if } m_{ij} < \infty \text{ and is odd, or} \\
\langle \sigma_j \sigma_i \rangle^{m_{ij}-1} \tau_i & \text{if } m_{ij} < \infty \text{ and is even,} \\
\tau_i \tau_j = \tau_j \tau_i & \text{if } m_{ij} = 2.
\end{cases}
\]

In arguments below we may regard a relation formally as an ordered pair of words. If \( X \) is a set of ordered pairs of words then \( X^{\Sigma} = \{(U, V) \mid (U, V) \text{ or } (V, U) \in X\} \).

The special case when \( I = \{1, \ldots, n\} \), \( m_{ij} = 3 \) when \( |i - j| = 1 \) and \( m_{ij} = 2 \) when \( |i - j| \geq 2 \) is the singular Artin monoid of type \( A_n \) and may be familiar to some readers as the singular braid monoid on \( n + 1 \) strings, denoted \( SB_{n+1} \), which was introduced by Baez and Birman in [2] and [4] respectively. We say \( M \) is right-angled if \( m_{ij} = \{2, \infty\} \) for \( i, j \in I \); and \( M = I_2(p) \) whenever \( I = \{1, 2\} \) and \( m_{12} = p \) for some \( p \geq 3 \).

If \( A \) and \( B \) are words in the above generators, we write \( A \approx B \) if \( A \) can be transformed into \( B \) by the use of the set of defining relations of \( SG_M \), and \( A = B \) if the two words are equal letter by letter.

We define the positive singular Artin monoid, denoted by \( SG_M^+ \), to be the monoid generated by \( S \cup T \) and the set \( \mathcal{R} \) of defining relations comprised of both \( \mathcal{R}_1 \) and \( \mathcal{R}_2 \) listed above. Where it does not cause confusion, elements of \( G_M, G_M^+, SG_M, \) and \( SG_M^+ \) may be referred to by words which represent them. If \( A \) and \( B \) are words in the generators \( S \) and \( T \), we write \( A \sim B \) if \( A \) can be transformed into \( B \) by the use of the set \( \mathcal{R} \) of defining relations. Theorems 1.1 (1) and (2) below are proved in [13] and [6] respectively:

**Theorem 1.1.**

(1) Let \( A, B \) be words over \( S \) such that \( A \approx B \). Then \( A \sim B \).

(2) Assume \( M \) is of finite type and let \( A, B \) be words over \( S \cup T \) such that \( A \approx B \). Then \( A \sim B \).

We denote by \( \ell(A) \) the length of any word \( A \) over \( S \cup T \). It is easy to see, by inspection of the defining relations, that \( SG_M^+ \) is homogeneous. We define the reduction property as in page 258 of [6]. By Lemma 15 of [6], cancellation and the reduction property hold in \( SG_M^+ \). By reduction we mean an application of the reduction property.

Let \( A \) and \( B \) be words in \( (S \cup T)^* \). We say \( A \) divides \( B \) or \( B \) is a multiple of \( A \) if there exists a word \( X \) in \( SG_M^+ \) such that \( B \sim AX \) in which case we write \( A \prec B \). We say \( A \) right divides \( B \), or \( B \) is a right multiple of \( A \), if there is a word \( X \) in \( SG_M^+ \) such that \( B \simXA \) in which case we write \( B \succ A \).

Let \( \Omega = \{A_1, A_2, \ldots, A_k\} \) be a set of words in \( (S \cup T)^* \). If \( \Omega \) has a common multiple then by Corollary 13 of [6], \( \Omega \) has a least common multiple (unique up to equivalence) which we denote by \( \text{lcm}(A_1, A_2, \ldots, A_k) \) or \( \text{lcm}(\Omega) \). By homogeneity,
lcm(\Omega) when it exists has minimal length. If \Omega has no common multiple then we write lcm(A_1, A_2, \ldots, A_k) = \infty.

Define a map \varphi from SG_M to the group algebra \mathbb{Z}[G_M] induced by
\[
\sigma_i^{\pm 1} \mapsto \sigma_i^{\pm 1}, \quad \tau_i \mapsto \sigma_i^2 + 1 \quad \text{for } i \in I.
\]

Then \varphi is easily verified to be a monoid homomorphism.

In Section 2, the Vassiliev homomorphism \eta is defined, we state Birman’s conjecture and show that \eta may be replaced (up to composition) with the simpler homomorphism \varphi. Moreover, for finite type \( M \), we show that if \varphi is injective on \( SG_M^+ \) then Birman’s conjecture is true for \( SG_M^+ \); this is followed by some observations regarding \varphi. In Section 3, we study the relationship between divisibility in \( SG_M^+ \) and the support of \varphi. In particular, we deduce that if \( U, V \in SG_M^+, C \in G_M^+ \) and \( \varphi(U) = \varphi(V) \) then \( C \) divides \( U \) if and only if \( C \) divides \( V \). Furthermore, we obtain that in \( SG_M^+ \), \varphi is injective on pairs of words for which a common multiple exists. In Section 4, we prove that \varphi preserves the Intermediate Property which holds in \( SG_M^+ \), namely, if \( \varphi(\tau_i U) = \varphi(\tau_i V) \) then \( m_{ij} = 2 \) or \( i = j \). From this it follows that \varphi is injective for some classes of monoids one of which includes singular Artin monoids of type \( I_2(p) \) generalising a result of East [7]. Finally, in Section 5, we discover that when \( M \) is of any (not necessarily finite) type, \( SG_M^+ \) injects into \( SG_M \) extending Theorem 1.1(2).

## 2 Birman’s conjecture

The map \( \eta \) from \( SB_{n+1} \) to the group algebra \( \mathbb{Z}[B_{n+1}] \) induced by
\[
\sigma_i^{\pm 1} \mapsto \sigma_i^{\pm 1}, \quad \tau_i \mapsto \sigma_i - \sigma_i^{-1}
\]
is easily verified to be a monoid homomorphism; \( \eta \) is sometimes referred to as the Vassiliev homomorphism [15] or desingularisation map [14]. In Remark 1 of [4], Birman conjectured that \( \eta \) is faithful, so that the singular braid monoid embeds into the group algebra of the braid group. In [14], Paris proved the truth of the conjecture. In Remark 25 of [6], Corran observed that Birman’s conjecture generalises to arbitrary Artin types since the Vassiliev homomorphism, \( \eta \), from any singular Artin monoid to the group algebra of the corresponding Artin group is well defined by the above rule. In effect, Birman’s conjecture was shown to hold when \( M = I_2(p) \) and \( M \) is right-angled in [7] and [10] respectively. Thus we may conjecture the following:

**Conjecture 1.** The Vassiliev homomorphism \( \eta : SG_M \longrightarrow \mathbb{Z}[G_M] \) is faithful, so that the singular Artin monoid embeds into the group algebra of the Artin group.

We write \( X = Y \) if \( X \) and \( Y \) are equal elements of \( \text{Im}(\eta) \) which should not cause confusion with our use of equality for words.

For every \( i \in I \) let \( \delta_i = \sigma_i \tau_i \), and put \( T' = \{ \delta_i \mid i \in I \} \). In [14], Paris observed that \( SG_M \) is generated as a monoid by \( S \cup S^{-1} \cup T' \) and admits a monoid presentation with
relations: \( \mathcal{R}_1 \); the free group relations; and the relations listed below, which are just \( \mathcal{R}_2 \) with \( \tau \) replaced by \( \delta \):

\[
\delta_i \sigma_i = \sigma_i \delta_i \quad \text{for all } i \in I,
\]

\[
\delta_i \langle \sigma_j \sigma_i \rangle^{m_{ij} - 1} \delta_j = \begin{cases} 
\langle \sigma_j \sigma_i \rangle^{m_{ij} - 1} \delta_j & \text{if } m_{ij} < \infty \text{ and is odd, or} \\
\langle \sigma_j \sigma_i \rangle^{m_{ij} - 1} \delta_i & \text{if } m_{ij} < \infty \text{ and is even,}
\end{cases}
\]

\[
\delta_i \delta_j = \delta_j \delta_i \quad \text{if } m_{ij} = 2.
\]

Thus the mapping \( \sigma^\pm_1 \mapsto \sigma^\pm_1 \), \( \tau_i \mapsto \delta_i \) for \( i \in I \) induces an isomorphism \( \mu : SG_M \to SG_M \), \( \sigma^\pm_1 \mapsto \sigma^\pm_1 \), \( \tau_i \mapsto \delta_i \). Furthermore, the Vassiliev homomorphism \( \eta : SG_M \to \mathbb{Z}[G_M] \) is determined by

\[
\eta(\sigma^\pm_1) = \sigma^\pm_1, \quad \eta(\delta_i) = \sigma_i^2 - 1 \quad \text{for } i \in I.
\]

The following Lemma 2.1 tells us that (up to composition) \( \eta \) may be replaced with the simpler homomorphism \( \varphi \). The proof of the lemma employs Lemma 1(1) of [11]; there it is proved only for type \( A_n \), but the argument is the same for any \( M \).

**Lemma 2.1.** \( \eta \) is injective precisely when \( \varphi \) is injective, and for any elements \( A, B \) of \( SG_M \),

\[
\varphi(A) = \varphi(B) \text{ if and only if } \eta(\mu(A)) = \eta(\mu(B)).
\]

**Proof:** Define the homomorphism, introduced in [11], \( \psi : SG_M \to \mathbb{Z}[G_M] \) by \( \psi(\sigma^\pm_1) = \sigma^\pm_1 \) and \( \psi(\tau_i) = \sigma_i + \sigma_i^{-1} \). Notice that for every \( i \in I \),

\[
\varphi(\sigma^\pm_1) = \sigma^\pm_1 \quad \text{and} \quad \varphi(\tau_i) = \psi(\delta_i),
\]

so that \( \varphi = \psi \mu \). Moreover, by Lemma 1(1) of [11],

\[
\eta(A) = \eta(B) \iff \psi(A) = \psi(B),
\]

and the result now follows readily. \( \square \)

Define the homomorphism \( \overline{\varphi} \) from \( SG_M^+ \) to \( \mathbb{Z}[G_M^+] \) by

\[
\overline{\varphi} : \sigma_i \mapsto \sigma_i, \quad \tau_i \mapsto \sigma_i^2 + 1 \quad i \in I.
\]

Thus we may conjecture the following:

**Conjecture 2.** \( \overline{\varphi} : SG_M^+ \to \mathbb{Z}[G_M^+] \) is injective.

Observation 1 below tells us that Conjecture 2 implies Conjecture 1 whenever \( M \) is of finite type.

Let \( U \) be any word over \( S \cup S^{-1} \cup T \) also regarded as an element of \( SG_M \). A *summand* of \( \varphi(U) \) is any word over \( S \cup S^{-1} \) obtained by replacing any given instance of \( \tau \) by \( \sigma^2 \) or \( 1 \). The *support* of \( \varphi(U) \) is the set of summands of \( \varphi(U) \). For example, in type \( A_2 \), \( \varphi(\tau_1 \sigma_2 \tau_1) \) has summands \( \sigma_1^2 \sigma_2 \sigma_1^2, \sigma_2 \sigma_1^2, \sigma_1^2 \sigma_2 \) and \( \sigma_2 \).

The ensuing result is a consequence of Theorem 1.1(1):
Corollary 2.2. Let \(U, V\) be words over \(S \cup T\), regarded alternatively as elements of \(SG_M\) and \(SG^+_M\), such that \(\varphi(U) = \varphi(V)\). Then \(\overline{\varphi}(U) = \overline{\varphi}(V)\).

Proof: Let \(F\) be any summand of \(\varphi(U)\). Then \(F\) is a word over \(S\), and there exists an element \(F'\) in the support of \(\varphi(V)\) such that \(F \approx F'\). Since \(F'\) is also over \(S\), Theorem 1.1(1) gives \(F \sim F'\). Hence, for every element \(F\) in the support of \(\varphi(U)\) there is a corresponding summand \(F'\) of \(\varphi(V)\) such that \(F \sim F'\). By regarding \(U\) and \(V\) also as elements of \(SG^+_M\), we infer that \(\overline{\varphi}(U) = \overline{\varphi}(V)\) as required. \(\square\)

Define a map \(\overline{+}\) from \((S \cup T)^*\) to \(S^*\) by \(\overline{+} : \sigma_i \mapsto \sigma_i\),\(, \tau_i \mapsto \tau_i^2\), for every \(i \in I\). Then it is easily verified that \(\overline{+}\) induces a homomorphism \(SG^+_M \to G^+_M\).

Remark 1. Whenever \(U, V\) are elements of \((S \cup T)^*\) with the same image under \(\overline{\varphi}\), \(U^+ \sim V^+\) as they both represent the unique monomial of maximal exponent sum of \(\overline{\varphi}(U) = \overline{\varphi}(V)\).

Let \(\Delta = \text{lcm}(\sigma_1, \ldots, \sigma_n)\). Theorem 5.6 of [5] tells us that \(\Delta\) exists precisely when \(M\) is of finite type, whilst by Section 5 of [6] the following holds:

Theorem 2.3. Let \(M\) be of finite type. For any word \(W\) over \(S \cup S^{-1} \cup T_1\), where \(T_1 \subseteq T\), there exists a word \(\overline{W}\) over \(S \cup T_1\) and integer \(p\) such that \(W \approx \Delta^p \overline{W}\).

Observation 1. Whenever \(M\) is of finite type, Conjecture 2 implies Conjecture 1.

Proof: Suppose \(M\) is of finite type, Conjecture 2 holds and that \(\eta(U) = \eta(V)\) for some words \(U\) and \(V\) in \((S \cup S^{-1} \cup T)^*\), where without causing confusion, we denote the equivalence class of a word by the word itself. By noting that \(\mu\) is an isomorphism, we deduce the existence of words \(U', V'\) such that \(U = \mu(U')\) and \(V = \mu(V')\); thus,

\[
\eta(\mu(U')) = \eta(U) = \eta(V) = \eta(\mu(V')),
\]

so by Lemma 2.1, \(\varphi(U') = \varphi(V')\). By Theorem 2.3, there exist integers \(p\) and \(q\), and words \(P, Q\) over \(S \cup T\) such that \(U' \approx \Delta^p P\) and \(V' \approx \Delta^q Q\). Hence there exist positive integers \(k_1, k_2\) and \(k\) such that

\[
\Delta^{k_1} U' \approx \Delta^{k_1} P \quad \text{and} \quad \Delta^{k_2} V' \approx \Delta^{k_2} Q.
\]

Thus (by recalling \(\Delta\) is over \(S\)),

\[
\varphi(\Delta^{k_1} P) = \varphi(\Delta^k U') = \varphi(\Delta^{k_2} V') = \varphi(\Delta^{k_2} Q).
\]

But \(k_1, k_2\) are positive integers and \(P, Q\) are over \(S \cup T\), so by Corollary 2.2, \(\overline{\varphi}(\Delta^{k_1} P) = \overline{\varphi}(\Delta^{k_2} Q)\), whence \(\Delta^{k_1} P \sim \Delta^{k_2} Q\) (since Conjecture 2 is assumed to be true). By (2.1), it follows that \(\Delta^k U' \approx \Delta^k V'\). By cancellation, \(U' \approx V'\), and hence

\[
U = \mu(U') \approx \mu(V') = V
\]

as required. \(\square\)
3 Common divisors and the support of \( \varphi \)

The following Lemma 3.1 tells us common “factors” of elements of \( SG_M \) with the same image under \( \varphi \) may be “cancelled out”. The proof of the lemma employs Lemmas 1(1) and 2 of [11]; there they are proved only for type \( A_n \), but the proofs proceed unmodifed for any \( M \).

**Lemma 3.1.** Let \( C, A \) and \( B \) be elements of \( SG_M \) such that \( \varphi(CA) = \varphi(CB) \). Then \( \varphi(A) = \varphi(B) \).

**Proof:** By Lemma 2.1,

\[
\eta(\mu(C)\mu(A)) = \eta(\mu(CA)) = \eta(\mu(CB)) = \eta(\mu(C)\mu(B)),
\]

whence \( \eta(\mu(A)) = \eta(\mu(B)) \) by Lemmas 1(1) and 2 of [11]. This implies that \( \varphi(A) = \varphi(B) \), again by Lemma 2.1, as required. \( \square \)

Let \( U \) be any word over \( S \cup T \), also regarded as an element of \( SG_M^+ \). We call \( U \) prime if there exists an \( s \in I \) such that \( \sigma_s \not\in U \). Let \( \mathcal{M}_U \) denote the summands of \( \varphi(U) \) that are prime. For example, in type \( A_2 \), \( \varphi(\tau_1\tau_2\tau_1) \) has summands \( \sigma_1^2\sigma_2\sigma_1^2, \sigma_1\sigma_2\sigma_1, \sigma_2\sigma_1, \) and a routine calculation shows

\[
\mathcal{M}_{\tau_1\tau_2\tau_1} = \{ \sigma_1^2\sigma_2, \sigma_2\sigma_1, \sigma_2 \}.
\]

**Lemma 3.2.** Let \( U, V \) be words in \( (S \cup T)^* \), also regarded as elements of \( SG_M^+ \), such that \( \varphi(U) = \varphi(V) \). For every summand \( F \) of \( \varphi(U) \) there is a corresponding summand \( G \) of \( \varphi(V) \) such that \( F \sim G \), and \( F \in \mathcal{M}_U \) if and only if \( G \in \mathcal{M}_V \).

**Proof:** Let \( F \) be any element in the support of \( \varphi(U) \). Then there exists an element \( G \) in the support of \( \varphi(V) \) such that \( F \sim G \). Furthermore, if \( F \in \mathcal{M}_U \), \( \sigma_i \not\in F \sim G \) for some \( i \in I \); this implies that \( G \) is prime and so must be an element of \( \mathcal{M}_V \). \( \square \)

Lemma 3.3 below is a preliminary result to Proposition 3.4.

**Lemma 3.3.** Let \( W \) be a word over \( S \cup T \), \( a \in I \), and suppose \( \sigma_a \not\in W \). Then there exists a word \( Z \in \mathcal{M}_W \) that is not divisible by \( \sigma_a \).

**Proof:** We prove the result by induction on \( \ell(W) \) which holds trivially if \( W \) is the empty word. So assume \( W \neq 1 \), and as an inductive hypothesis, that the result holds for words of length \( \ell(W) \). Then \( W = \alpha_rW_1 \) for some word \( W_1 \), \( r \in I \) and \( \alpha = \sigma \) or \( \tau \).

Suppose first that \( r = a \). Then \( \sigma_a \not\in W = \alpha_aW_1 \) gives \( \alpha = \tau \), and since \( \sigma_a \) and \( \tau_a \) commute, \( \sigma_a \not\in W_1 \). By the inductive hypothesis, we infer the existence of a word \( Y \in \mathcal{M}_{W_1} \) such that \( \sigma_a \not\in Y \). Notice that since \( W = \tau_aW_1 \), \( Y \) is also an element of \( \mathcal{M}_W \). Putting \( Z = Y \) thus gives the desired result.

Assume then that \( m_{ra} \geq 2 \). There exists a word \( W_2 \) and a largest integer \( q \) such that

\[
W = \alpha_rW_1 \sim \alpha_r(\sigma_a\sigma_r)^qW_2 \quad \text{and} \quad \sigma_c \not\in W_2, \quad (3.1)
\]
where \( c = a \) if \( q \) is even and \( c = r \) if \( q \) is odd. By the inductive hypothesis, we deduce the existence of a word \( Y \) in \( M_{W_2} \) such that \( \sigma_c \not\preceq Y \). Observe that \( \sigma_a \prec \alpha_r (\sigma_a \sigma_r)^{m_{ra} - 1} \) whenever \( m_{ra} < \infty \) and that \( \sigma_a \not\prec \alpha_r W_1 = W \); this tells us that \( 0 \leq q \leq m_{ra} - 2 \). Put \( Y' = \lambda_r (\sigma_a \sigma_r)^q Y \) where \( \lambda_r = \sigma_r^2 \) if \( \alpha = \tau \) and \( \lambda_r = \sigma_r \) if \( \alpha = \sigma \). Then \( Y' \) is clearly a summand of \( \overline{\varphi}(\alpha_r(\sigma_a \sigma_r)^q W_2) \). Furthermore, if \( \sigma_a \) were to divide \( Y' \), reduction would show that \( \sigma_c \) divides \( Y \) which would contradict that \( \sigma_c \not\preceq Y \). This implies \( \sigma_a \not\preceq Y' \), so \( Y' \in M_{(\sigma_a \sigma_r)^q W_2} \). Certainly \( \overline{\varphi}(W) = \overline{\varphi}(\alpha_r(\sigma_a \sigma_r)^q W_2) \), where by (3.1), \( W \) and \( \alpha_r(\sigma_a \sigma_r)^q W_2 \) are regarded as (the same) elements of \( SG^+_M \). By Lemma 3.2, we deduce that \( \sigma_a \not\preceq Y' \sim Z \) for some word \( Z \in M_W \) as required, and the result follows by induction.

**Proposition 3.4.** Let \( U, V \) be words in \( (S \cup T)^* \), also regarded as elements of \( SG_M \), such that \( \varphi(U) = \varphi(V) \), and let \( C \) be any word over \( S \). Then \( C \) divides \( U \) if and only if \( C \) divides \( V \).

**Proof:** We first prove the “only if” part of the result. By Corollary 2.2, \( \overline{\varphi}(U) = \overline{\varphi}(V) \). Let \( s \in I \) and suppose that \( \sigma_s \prec U \) yet \( \sigma_s \not\prec V \). Then \( U \sim \sigma_s U_1 \) for some word \( U_1 \), and by Lemma 3.3, there exists a word \( X \in M_V \) such that \( \sigma_s \not\prec X \). By Lemma 3.2, we infer the existence of a word \( Y \in M_{s U_1} \) such that \( X \sim Y \). But \( Y \) is a summand of \( \overline{\varphi}(\sigma_s U_1) \), whence \( Y = \sigma_a Y' \) for some word \( Y' \) in the support of \( \overline{\varphi}(U_1) \). Thus \( \sigma_s \not\prec X \sim Y = \sigma_s Y' \), a contradiction; so \( \sigma_s \prec U \) only if \( \sigma_s \prec V \). This proves the result for \( \ell(C) = 1 \), and by noting that it holds trivially for \( \ell(C) = 0 \), starts an induction. So assume that \( C \) divides \( U \) and \( \ell(C) \geq 2 \). Then there exists a letter \( a \) in \( S \) and non-empty word \( C_1 \) over \( S \) such that \( C = C_1 a \) and \( U \sim C_1 a U_1 \) for some word \( U_1 \) over \( S \cup T \). By the inductive hypothesis, we derive the existence of a word \( V_1 \) over \( S \cup T \), such that \( V \sim C_1 V_1 \). Thus, \( \varphi(C_1 a U_1) = \varphi(U) = \varphi(V) = \varphi(C_1 V_1) \), whence \( \varphi(a U_1) = \varphi(V_1) \) by Lemma 3.1, so that \( a \) divides \( V_1 \). Hence \( C = C_1 a \) also divides \( V \sim C_1 V_1 \) as required, and the result for any \( \ell(C) \) follows by induction. Swapping the roles of \( U \) and \( V \) in the preceding argument gives the “if” part of the claim.

The proof of the following result proceeds identically to that of Proposition 2.1 of [1]:

**Lemma 3.5.** Let \( j \in I \) and \( W \) any word in \( (S \cup T)^* \) such that \( \text{lcm}(\tau_j, W) \) exists. Then \( \sigma_j \) divides \( W^+ \) precisely when either \( \sigma_j \) or \( \tau_j \) divides \( W \).

**Lemma 3.6.** Let \( U, V \) be words in \( (S \cup T)^* \) also regarded as elements of \( SG_M^+ \) such that \( \overline{\varphi}(U) = \overline{\varphi}(V) \). Suppose \( \tau_j \) divides \( U \) and \( \text{lcm}(\tau_j, V) \) exists. Then either \( \sigma_j \) or \( \tau_j \) is a common divisor of \( U \) and \( V \).

**Proof:** We have \( \sigma_j \prec U^+ \sim V^+ \), so by Lemma 3.5, we infer that \( V \sim \alpha_j V' \) for some word \( V' \) over \( S \cup T \), \( \alpha = \sigma \) or \( \tau \). If \( \alpha = \tau \), it is evident that \( \tau_j \) is a common divisor of \( U \) and \( V \); whilst if \( \alpha = \sigma \) then by Proposition 3.4 we see that \( \sigma_j \) also divides \( U \) and hence is a common divisor of \( U \) and \( V \).
Define monoid homomorphisms $\epsilon$ and $\mathcal{N}$ from $SG_M$ to $(\mathbb{Z}, +)$ by

$$
\epsilon : \sigma_i^{+1} \mapsto 1, \; \tau_i \mapsto 0, \; \mathcal{N} : \sigma_i^{+1} \mapsto 0, \; \tau_i \mapsto 1 \quad \text{for } i \in I.
$$

So $\epsilon(A)$ is the exponent sum of sigmas and $\mathcal{N}(A)$ counts the number of singularities (taus) of any word $A$ in $SG_M$.

**Lemma 3.7.** Let $U$ and $V$ be words over $S \cup T$ regarded as elements of $SG_M$, and suppose $\varphi(U) = \varphi(V)$. Then $\ell(U) = \ell(V)$ and $\mathcal{N}(U) = \mathcal{N}(V)$.

**Proof:** By Corollary 2.2, $\overline{\varphi}(U) = \overline{\varphi}(V)$, whence $U^+ \sim V^+$. Thus

$$
\epsilon(U) + 2\mathcal{N}(U) = \ell(U^+) = \ell(V^+) = \epsilon(V) + 2\mathcal{N}(V).
$$

(3.2)

Notice that for every word $A$ over $S \cup T$ there is a unique monomial, represented by $A^-$, obtained by replacing each $\tau_i$ by 1, in the support of $\varphi(A)$ with minimal exponent sum $\epsilon(A)$. Then since $\overline{\varphi}(U) = \overline{\varphi}(V)$, it follows that $U^- \sim V^-$ which gives

$$
\epsilon(U) = \ell(U^-) = \ell(V^-) = \epsilon(V).
$$

(3.3)

By (3.2) and (3.3), we infer that $\mathcal{N}(U) = \mathcal{N}(V)$, whence

$$
\ell(U) = \epsilon(U) + \mathcal{N}(U) = \epsilon(V) + \mathcal{N}(V) = \ell(V)
$$

as required. \qed

**Theorem 3.8.** Over $S \cup T$, $\varphi$ is injective on pairs of words for which a common multiple exists.

**Proof:** Let $U, V$ be words in $(S \cup T)^*$ such that $\varphi(U) = \varphi(V)$ and $\text{lcm}(U, V)$ exists. We show $U \sim V$ by induction on $\ell(U)$. By Corollary 2.2, $\overline{\varphi}(U) = \overline{\varphi}(V)$. If $U$ is the empty word, there is nothing to show, so suppose $\ell(U) \geq 1$. If $\tau_j \prec U$ then either $\sigma_j$ or $\tau_j$ is a common divisor of $U$ and $V$ by Lemma 3.6; whilst if $\sigma_j \prec U$ then $\sigma_j$ also divides $V$ by Proposition 3.4. In either case we see that $U \sim \alpha_j U'$ and $V \sim \alpha_j V'$ for some words $U', V'$ and $\alpha = \sigma$ or $\tau$. By Lemma 3.1, this implies that $\varphi(U') = \varphi(V')$, and since $\text{lcm}(\alpha_j U', \alpha_j V') \sim \text{lcm}(U, V) \neq \infty$, we obtain that $\text{lcm}(U', V')$ also exists. By the inductive hypothesis, we deduce that $U' \sim V'$, whence $U \sim \alpha_j U' \sim \alpha_j V' \sim V$ as required, and the result now follows by induction. \qed

## 4 The Intermediate Lemma

In this section we prove that the Intermediate Property – discovered in Intermediate Lemma of [6] and expressed in Lemma 4.1 below – is preserved under $\varphi$. As a corollary we deduce that $\varphi$ is injective for a class of monoids which include singular Artin monoids of type $I_2(p)$. 
Lemma 4.1. Let $U, V$ be words in $(S \cup T)^*$ such that $\tau_i U \sim \tau_j V$. Then $i = j$ or $m_{ij} = 2$.

The ensuing Lemma 4.2 is a preliminary result to Theorem 4.3 below.

Lemma 4.2. Let $F, G \in (S \cup T)^*$, $i, j \in I$, and suppose $\sigma_j \not< \sigma_i^2 F \sim G$. Then either $\sigma_i \not< \sigma_j^2 G$ or $m_{ij} = 2$.

Proof: Cleary $m_{ij} \geq 2$ since $\sigma_j \not< \sigma_i^2 F$. So suppose $m_{ij} \geq 3$ yet $\sigma_j^2 G \sim \sigma_i H$ for some word $H$. By reduction, $\sigma_i \not< \sigma_j G$, which gives by another reduction, $G \sim \langle \sigma_i \sigma_j \rangle^{m_{ij}-1} G'$ for some word $G'$. Thus,

$$\sigma_j^2 G \sim G \sim \langle \sigma_i \sigma_j \rangle^{m_{ij}-1} G' = \sigma_i \langle \sigma_j \sigma_i \rangle^{m_{ij}-2} G'.$$

By noting that $m_{ij} \geq 3$, cancellation applied to the preceding equivalence shows that $\sigma_j \not< \sigma_i F$, so by reduction,

$$F \sim \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} F_1 \tag{4.1}$$

for some word $F_1$. Put $c = j$ if $m_{ij}$ is odd and $c = i$ if $m_{ij}$ is even. Then

$$\sigma_j \not< \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} \sigma_i^2 F_1 \sim \sigma_i^2 \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} F_1 \sim \sigma_i^2 F$$

by (4.1), and this contradicts the hypothesis. \hfill $\square$

Theorem 4.3. Let $U = \tau_i U_1$, $V = \tau_j V_1$ be words in $(S \cup T)^*$, also regarded as elements of $SG_M$, such that $\varphi(U) = \varphi(V)$. Then $i = j$ or $m_{ij} = 2$.

Proof: Suppose $U = \tau_i U_1$ and $V = \tau_j V_1$ provide a counterexample. That is, $\varphi(\tau_i U_1) = \varphi(\tau_j V_1)$ but $m_{ij} \geq 3$. Suppose further that this counterexample is minimal with respect to $\ell(U)$ which by Lemma 3.7 is equal to $\ell(V)$. Then $\ell(U) \geq 2$, since $\ell(U) = \ell(V) = 1$ gives $\sigma_i^2 + 1 = \varphi(U) = \varphi(V) = \sigma_j^2 + 1$ and shows $i = j$. We first prove that $V = \tau_j V_1$ is not divisible by $\sigma_j$. Suppose, by way of contradiction, that it is. Reduction yields a word $P$ such that $V_1 \sim \sigma_j P$, and by recalling that $\varphi(U) = \varphi(V)$, Proposition 3.4 implies that $\sigma_j$ also divides $U = \tau_i U_1$, so by reduction again, $U_1 \sim \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} Q$ for some word $Q$. Put $C = \langle \sigma_j \sigma_i \rangle^{m_{ij}-1}$. Then

$$U = \tau_i U_1 \sim \tau_i C Q \sim C \tau_d Q \tag{4.2}$$

where $d = j$ if $m_{ij}$ is odd and $d = i$ if $m_{ij}$ is even. Since $C$ is over $S$, Proposition 3.4 may be applied and this shows that

$$C = \sigma_j \langle \sigma_i \sigma_j \rangle^{m_{ij}-2} \sim V \sim \tau_j \sigma_j P \sim \sigma_j \tau_j P.$$

By recalling that $m_{ij} \geq 3$, we deduce (after cancelling) that $\sigma_i \not< \tau_j P$, so by reduction,

$$P \sim \langle \sigma_i \sigma_j \rangle^{m_{ij}-1} P'$$

for some word $P'$ over $S \cup T$. Thus,

$$V \sim \sigma_j \tau_j P \sim \sigma_j \tau_j \langle \sigma_i \sigma_j \rangle^{m_{ij}-1} P' \sim \sigma_j \langle \sigma_i \sigma_j \rangle^{m_{ij}-1} \tau_c P' \sim \langle \sigma_j \sigma_i \rangle^{m_{ij}-1} \sigma_d \tau_c P' = C \sigma_d \tau_c P'.$$
where \( m_{cd} \geq 3 \) (since \( \{c, d\} = \{i, j\} \)). When combined with (4.2), this implies

\[
\varphi(C\sigma_d\tau_c P^\prime) = \varphi(V) = \varphi(U) = \varphi(C\tau_c Q),
\]
so by Lemma 3.1, \( \varphi(\sigma_d\tau_c P^\prime) = \varphi(\tau_c Q) \). By invoking Proposition 3.4 again, we infer that \( \sigma_d \) divides \( \tau_c Q \), and by reduction, this yields a word \( Q' \) such that \( Q \sim \sigma_d Q' \). Hence

\[
\varphi(\sigma_d\tau_c P^\prime) = \varphi(\tau_c Q) = \varphi(\tau_c \sigma_d Q') = \varphi(\sigma_d \tau_c Q')
\]
so by Lemma 3.1 again, \( \varphi(\tau_c P^\prime) = \varphi(\tau_c Q') \), and \( m_{cd} \geq 3 \). Since \( \ell(V) > \ell(\tau_c P^\prime) \), this contradicts the minimality of \( \ell(U) = \ell(V) \). Thus \( \sigma_j \not\sim V \), and so by a final application of Proposition 3.4, we deduce that \( \sigma_j \not\sim U \). Observe that \( \overline{\tau}(U) = \overline{\tau}(V) \) by Corollary 2.2.

Now let \( F \) be an element of \( \mathcal{M}_U \) such that \( \sigma_j \not\sim F \), the existence of which is assured by Lemma 3.3. Assume further that \( \ell(F) \) is maximal; that is,

\[
\text{if } \sigma_j \not\sim F' \text{ and } F' \in \mathcal{M}_U \text{ then } \ell(F') \leq \ell(F). \tag{4.3}
\]

By noting that \( U = \tau_i U_1 \), we infer the existence of a summand \( F_1 \) of \( \overline{\tau}(U_1) \) such that either \( F = \sigma_i^2 F_1 \) or \( F = F_1 \). Suppose that \( F = F_1 \), and put \( F' = \sigma_i^2 F_1 \). Then \( F' \) is clearly an element in the support of \( \overline{\tau}(U) \) and \( \ell(F') > \ell(F) \). By (4.3), we deduce that \( \sigma_j \not\sim \sigma_i^2 F_1 \). Since \( m_{ij} \geq 3 \), the reduction property implies that \( \sigma_j \) divides \( F_1 = F \), which contradicts that \( \sigma_j \not\sim F \). This shows that \( F = \sigma_i^2 F_1 \). By Lemma 3.2, there exists a word \( G \in \mathcal{M}_V \) such that \( F \sim G \). By recalling \( V = \tau_j V_1 \), we deduce that either \( G = \sigma_j^2 G_1 \) or \( G = G_1 \), for some word \( G_1 \) in the support of \( \overline{\tau}(V_1) \). The first possibility must be excluded since \( \sigma_j \not\sim F \sim G \), so \( G = G_1 \). Thus,

\[
\sigma_j \not\sim \sigma_i^2 F_1 = F \sim G = G_1 \text{ and } G_1 \in \mathcal{M}_V. \tag{4.4}
\]

Put \( G' = \sigma_j^2 G_1 \), and note that it is a summand of \( \overline{\tau}(V) = \overline{\tau}(\tau_j V_1) \). By Lemma 4.2, \( \sigma_i \not\sim G' \), so \( G' \) is an element of \( \mathcal{M}_V \). By Lemma 3.2, \( G' \sim F' \) for some word \( F' \in \mathcal{M}_U \), and by (4.4), \( \ell(G') = 2 + \ell(G_1) = 2 + \ell(F) \). Hence

\[
\sigma_i \not\sim \sigma_j^2 G_1 = G' \sim F', \quad F' \in \mathcal{M}_U \quad \text{and} \quad \ell(F') = 2 + \ell(F). \tag{4.5}
\]

Since \( F' \) is a summand of \( \overline{\tau}(U) = \overline{\tau}(\tau_i U_1) \), we infer the existence of a word \( F_2 \) in the support of \( U_1 \) such that either \( F' = \sigma_i^2 F_2 \) or \( F' = F_2 \). As the first possibility evidently contradicts (4.5), we deduce that \( F' = F_2 \). Put \( F'' = \sigma_i^2 F_2 \), and note that it is a summand of \( \overline{\tau}(\tau_i U_1) = \overline{\tau}(U) \). By Lemma 4.2 and (4.5), \( \sigma_j \not\sim \sigma_i^2 F' = F'' \), whence the latter word is an element of \( \mathcal{M}_U \). When combined with (4.3), this implies that \( \ell(F'') \leq \ell(F) \). However, \( F'' = \sigma_i^2 F_2 = \sigma_i^2 F' \), so by (4.5),

\[
\ell(F'') = 2 + \ell(F') = 4 + \ell(F),
\]
a contradiction. \( \square \)

Fix \( i, j \in I \) such that \( m_{ij} \geq 3 \). Let \( T_{ij} \) and \( T_{ij}^+ \) denote the sets of equivalence classes of words in \( (S \cup S^{-1} \cup \{\tau_i, \tau_j\})^* \) and \( (S \cup \{\tau_i, \tau_j\})^* \) under \( \approx \) and \( \sim \) respectively. Then \( T_{ij} \) and \( T_{ij}^+ \) are both submonoids of \( SG_M \) and \( SG_M^+ \) respectively.
Corollary 4.4. The restriction of \( \varphi \) from \( T_{ij}^+ \) to the group algebra \( \mathbb{Z}[G_M^+] \) is injective. In particular, the desingularisation map \( \eta : SG_{I_2(p)} \to \mathbb{Z}[G_{I_2(p)}] \) is injective.

Proof: Suppose that \( U, V \) in \( (S \cup \{\tau_i, \tau_j\})^* \) provide a counterexample. That is, assume that \( U \neq V \) yet \( \varphi(U) = \varphi(V) \). Note that trivially \( \varphi(U) = \varphi(V) \) where \( U, V \) are also regarded as elements of \( SG_M \). Suppose further that this counterexample is minimal with respect to \( \ell(U) \) which, by Lemma 3.7, is the same as \( \ell(V) \). Certainly \( \ell(U) \geq 2 \). If \( U \sim CU' \), \( V \sim CV' \) for some non-empty word \( C \) then \( \varphi(U') = \varphi(V') \) by Lemma 3.1 and Corollary 2.2, \( U' \neq V' \), \( \ell(U') < \ell(U) \), and hence the minimality of \( \ell(U) \) is contradicted. Thus \( U, V \) have no non-trivial divisor from which we infer, by Proposition 3.4, that \( U \) and \( V \) are not divisible by any generator from \( S \). This tells us that \( U = \tau_r U_1 \) and \( V = \tau_s V_1 \) for some words \( U_1, V_1 \) in \( T_{ij}^+ \) and generators \( \tau_r, \tau_s \in \{\tau_i, \tau_j\} \). Since \( m_{ij} \geq 3 \), we deduce from Theorem 4.3 that \( r = s \), so \( \tau_r \) is a common divisor of \( U \) and \( V \). This is a contradiction since \( U \) and \( V \) have no non-trivial common divisor. By noting \( I_2(p) \) is of finite type, the last statement of the result follows by Observation 1, since \( T_{12}^+ = SG_{I_2(p)}^+ \). \( \square \)

5 Embedding \( SG_M^+ \) into \( SG_M \)

Recall that in [13], Paris discovered that all positive monoids inject into their groups (Theorem 1.1(1)), whilst in [6], Corran showed that \( SG_M^+ \) embeds into \( SG_M \) whenever \( M \) is of finite type (Theorem 1.1(2)). It would thus seem natural to ask whether the latter result can be extended to all singular Artin monoids, which is what we derive in this section (Theorem 5.3).

Let \( U, V \) be words over \( S \cup S^{-1} \cup T \). We say \( U \) and \( V \) differ by an elementary transformation if there are words \( X \) and \( Y \) and a relation \( (\rho_1, \rho_2) \in \mathcal{R}^2 \) such that \( V = X \rho_1 Y \) and \( U = X \rho_2 Y \). We say that a word \( V \) is obtained from \( U \) by a trivial insertion if there are words \( X, Y \) and a letter \( a \in S \cup S^{-1} \) such that \( U = XY \) and \( V = Xaa^{-1}Y \). In this case we also say that \( U \) is obtained from \( V \) by a trivial deletion.

The reader is reminded of the monoid homomorphism \( \mathcal{N} \) from \( SG_M \) to \( (\mathbb{Z}, +) \) defined by \( \sigma_i^{\pm 1} \mapsto 0, \tau_i \mapsto 1 \). Thus \( \mathcal{N} \) counts the number of taus in any given word. Now let \( W \) be a word over \( S \cup S^{-1} \cup T \), and suppose \( \mathcal{N}(W) = k \geq 1 \). Then there are words \( W_i \) over \( S \cup S^{-1} \) and generators \( \tau_{a_i} \in T \) such that

\[
W = W_0 \tau_{a_1} W_1 \tau_{a_2} W_2 \ldots W_{k-1} \tau_{a_k} W_k.
\]

For \( r = 1, \ldots, k \), let

\[
\theta_r(W) = W_0 \tau_{a_1} W_1 \ldots W_{r-2} \tau_{a_{r-1}} W_{r-1} \sigma_a W_r \tau_{a_{r+1}} W_{r+1} \ldots W_{k-1} \tau_{a_k} W_k,
\]

hence \( \theta_r \) reduces the number of taus of \( W \) by 1.

Lemma 5.1. Let \( W, V \) be equivalent words over \( S \cup S^{-1} \cup T \), and suppose \( \mathcal{N}(W) = k \geq 1 \). Then for every \( r \in \{1, \ldots, k\} \) there exists \( s \in \{1, \ldots, k\} \) such that \( \theta_r(W) \approx \theta_s(V) \).
Proof: Let \( r \) be any integer such that \( 1 \leq r \leq k \). Since \( W \approx V \) there is a sequence \( Z_1, \ldots, Z_t \) of words over \( S \cup S^{-1} \cup T \) such that \( W = Z_1 \approx Z_2 \approx \ldots \approx Z_t = V \) and \( Z_{t+1} \) is obtained from \( Z_t \) by an elementary transformation or by a trivial insertion or deletion. If \( t = 1 \), the result is trivial and hence starts an induction. So suppose \( t \) is at least 2. If \( Z_1 \) and \( Z_2 \) differ by an elementary transformation, inspection of \( \mathcal{R} \) gives \( \theta_r(Z_1) \approx \theta_r(Z_2) \) if the relation involves any \( \sigma \) and either \( \theta_r(Z_1) \approx \theta_{r+1}(Z_2) \), \( \theta_r(Z_2) \) or \( \theta_{r-1}(Z_2) \) if the relation involves commuting \( \tau \)'s; whilst if \( Z_2 \) is obtained from \( Z_1 \) by a trivial insertion or deletion, it is evident that \( \theta_r(Z_1) \approx \theta_r(Z_2) \). By the inductive hypothesis, we deduce that \( \theta_r(Z_2) \approx \theta_s(Z_t) \) for some \( s \in \{1, \ldots, k\} \), whence

\[
\theta_r(W) = \theta_r(Z_1) \approx \theta_r(Z_2) \approx \theta_s(Z_t) = \theta_s(V)
\]
as required. The result now follows by induction. \( \square \)

In [8], it was shown that the singular braid monoid on \( n+1 \) strings (that is, the singular Artin monoid of type \( A_n \)) can be embedded in a group. The group constructed by the authors relies heavily on the geometry of singular braids in space; more specifically, it has a geometric interpretation as singular braids with two types of (cancelling) singularities. By employing purely algebraic methods, Paris [14] gave another proof of the fact that singular braid monoids inject into groups. This was shown (chronologically and with completely different proofs) in [3], [12] and [10]. An evident corollary of the fact that \( SG_M \) injects into a group is that left and right cancellation hold in \( SG_M \); namely,

**Proposition 5.2.** Let \( C, W, V \) be words over \( S \cup S^{-1} \cup T \) such that either \( CW \approx CV \) or \( WC \approx VC \). Then \( W \approx V \).  \(^1\)

**Theorem 5.3.** Let \( W, V \) be words over \( S \cup T \) such that \( W \approx V \). Then \( W \sim V \).

**Proof:** Put \( k = N(W) \). If \( k = 0 \), the result follows by Theorem 1.1(1) and starts an induction. So assume \( k \geq 1 \). Certainly \( \varphi(W) = \varphi(V) \). Suppose \( \text{lcm}(W, V) = \infty \). But \( W \sim CW', V \sim CV' \) for some words \( C, W' \) and \( V' \) over \( S \cup T \), such that \( W', V' \) have no (non-trivial) common divisor. Thus \( \text{lcm}(W', V') = \infty \) and

\[
CW' \sim W \approx V \sim CV',
\]

so by Proposition 5.2, \( W' \approx V' \); this implies \( \varphi(W') = \varphi(V') \), and by Corollary 2.2, \( \overline{\varphi}(W') = \overline{\varphi}(V') \). Since \( W' \) and \( V' \) have no common divisor, Proposition 3.4 yields

\[
W' = \tau_i W'' \quad \text{and} \quad V' = \tau_j V''
\]

for some distinct \( i, j \in J \) and words \( W'', V'' \) over \( S \cup T \). By Theorem 4.3, \( m_{ij} = 2 \) (since \( i \neq j \)), whilst Lemma 3.6 gives

\[
\text{lcm}(\tau_j, W') = \infty.
\]  \(^{(5.1)}\)

\(^1\)The reader is referred to Proposition 5.1 of [9] for a (geometric) demonstration of this result, when \( M = A_n \), without invoking the embedding of \((SB_{n+1} =) SG_{A_n} \) into a group.
But $\tau_i W'' = W' \approx V' = \tau_j V''$, so by Lemma 5.1,

$$\sigma_i W'' = \theta_1(\tau_i W'') \approx \theta_r(\tau_j V'') \quad \text{for some } r \in \{1, \ldots, k\}.$$

The inductive hypothesis now tells us that $\sigma_i W'' \sim \theta_r(\tau_j V'')$, whence

$$\sigma_i W'' \sim \begin{cases} 
\sigma_j V'' & \text{if } r = 1, \\
\tau_j \theta_{r-1}(V'') & \text{if } 2 \leq r \leq k.
\end{cases}$$

If $\sigma_i W'' \sim \sigma_j V''$, reduction implies that $\sigma_j \prec W''$, so by noting that $m_{ij} = 2$, $\tau_i$ and $\sigma_j$ commute, whence $\sigma_j \prec \tau_i W'' = W''$; Proposition 3.4 can now be invoked to deduce that $W'$ and $V'$ have a non-trivial common divisor, a contradiction. Hence $\sigma_i W'' \sim \tau_j \theta_{r-1}(V'')$. By recalling that $m_{ij} = 2$, another application of reduction gives $\tau_j \prec W''$, so that $\tau_j \prec \tau_i W'' = W''$; this clearly contradicts (5.1). Thus lcm$(W, V)$ exists, so $W \sim V$ by Theorem 3.8, and the result follows by induction.

References


