Fautastic reference: Iwahori-Matsumolo, Publ. IHES, 1965.

Assume most our root system is irreducible. A.  $\Leftrightarrow$  A.  $\Leftrightarrow$  A. Weight lattice:  $\mathcal{X} = \{ A \in \mathcal{Y}^* \mid \langle A, \kappa' \rangle \in \mathbb{Z} \mid \forall \alpha \}$ .

For  $x \in \overline{\Phi}$  have reflection  $S_{\alpha}(N) := N - \langle \alpha^{\vee}, N \rangle \times .$  Gr  $A = \emptyset$  Remark: We can and do equip  $b^{*}$  with a Euclidean when product s.t.  $S_{\alpha}$  is orthogonal.  $W_{\underline{\Phi}} = \langle S_{\alpha} \mid \alpha \in \overline{\Phi} \rangle = \langle S_{\alpha} \mid \alpha \in \Sigma \rangle .$  "huite Weyl groups"

 $W_e = W_{\xi} \times \mathcal{X}$  "extended alline Given  $A \in \mathcal{X}$  with  $t_A$  for branslation by A. Weyl graps"

For  $\alpha \in \Phi_+$ ,  $m \in \mathbb{Z}$  consider  $H_{\alpha,m} := \{ \Delta \mid (\Delta, \alpha^{\vee}) = m \}$ .

Sx, m (1) := 1 - (x, 1) x + m 20 x.

Note that so, is a reflection and lixes Ho, in.

This characterises Sa, m uniquely.

Also Sor, m = tma o Sox = Sox o tma . (\*)

Lemma: W is generated by Sa, m.

Proof: Clear hom (\*). because 20 is generated by

Lunna: W is a normal subgroup of We.

Proof: Clearly to (Harm) = Harma (1, m). Hence

ta sx, m ta = sx, m+ (x, a).

 $\partial \mathcal{C} = \text{Set of relieding hyperplanes } H_{\sigma,m}, \propto \in \Phi, m \in \mathbb{Z}$ .

We, Wack on 76 60.

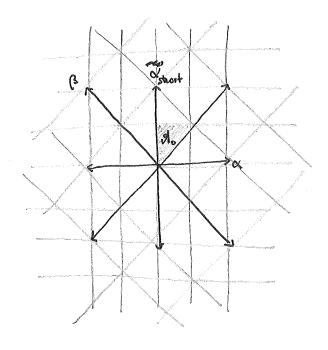
~> We, W aut on 1/8 \ HeTB

Connected components are called alcores.

while the coroot

 $A_0 = \{ N | \langle N, \kappa^{V} \rangle > 0 \ \forall \alpha \in \Sigma, \langle N, \langle \widetilde{\alpha}_{ab} \rangle^{V} \rangle < 1 \}$  lundamental alcove.

## Example:



Suppose some Hx, in intersects A. We can hind x+ B= 2. Choose vo in Mis uitersection.

Now:

 $M \leq \langle \alpha^{\nu}, \nu \rangle + \langle \beta^{\nu}, \nu \rangle = \langle \alpha^{\nu}, \nu \rangle < 1$ 

=> contradiction.

because or, or are sums of sniple rooks and  $v \in A_0$ . let S:= {Sα | α ∈ Δ} U {@S α short, 1} be the Rellutions in the walls of its.

Clauin:

DOG Q Lindow Sufak Cloubay . Con A Calling on Cy A. -> b\*/ (s>.

If ve its men we're done, otherwise Choose VE &\*. a higgeletane tepacition whose relluling exists mere

hyperplane seperales A. and

If pg devotes a point in me interior of As Men

11 s(v)-311 < 11v-311.

The set of Worloods of v is discrete,

hence 11 w·(v) -8 11 obtains a minimum.

This point must lie in do.

Now we'red done.

 $W = \langle S \rangle$ 

NOW PRINTED NOTES.

ye gely spigglest braughtively suggest; Courses of Disquelling, if Whites

· Two different interpretations of length builtion, simply mausitire.

$$||v - \rho|| > ||s(v) - \rho||.$$

Because  $W_S$  is discrete, there are finitely many points in the  $W_S$  orbit of  $\nu$  which are of distance at most  $||\nu - \rho||$  from  $\rho$ . Hence, using reflections from  $W_S$  we can keep reducing the distance from  $\rho$  to  $\nu$  until this is no longer possible, i.e. until  $\nu \in \Delta$ .

**Lemma 2.2.**  $W = W_S$ , i.e. W is generated by S.

*Proof.* Because W is generated by the reflections it contains, it is enough to show that any reflection in W belongs to  $W_S$ . To this end, fix  $\alpha \in \Phi$  and let r denote the corresponding affine reflection. Choose an alcove  $A \in \mathscr{A}$  such that  $F := A \cap \alpha$  is of dimension one less than V. By the previous lemma, there is an element  $w \in W_S$  such that  $wA = \Delta$ . Let  $s \in S$  be the reflection in the wall  $wF \subset wA = \Delta$ . Then

$$w^{-1}sw = r$$
.

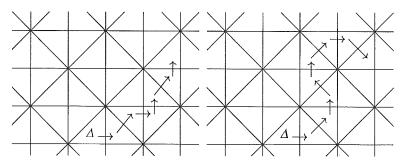
(Indeed, the left hand side is a reflection which fixes F and hence  $\alpha$ , and hence must be r.)

## 2.3 Expressions and strolls

Fix an affine reflection group W acting on V, together with a choice of fundamental alcove  $\Delta$  as above. Let S denote the set of reflections in the walls of  $\Delta$ . An expression for x is a word  $\underline{x} = (s_1, s_2, \ldots, s_m)$  in S such that  $x = s_1 s_2 \ldots s_m$ . The length  $\ell(\underline{x})$  of an expression is its length as a word. An expression for x is reduced if it is of minimal length amongst all possible expressions for x. The length  $\ell(x)$  is the length of a reduced expression.

A stroll is a sequence  $\underline{A} := (A_0, A_1, \dots, A_k)$  of elements of  $\overline{\mathscr{A}}$  such that  $A_0 = \Delta$  and  $A_{i-1}$  and  $A_i$  share a codimension 1 face  $F_i$  for all  $1 \le i \le k$ . We think of a stroll as a path in V beginning in  $\Delta$  and only passing through codimension 1 parts of the hyperplane arrangement  $\Phi$  (see the examples below). The  $length \ \ell(\underline{A})$  is the number of hyperplanes crossed by the path (i.e. if  $\underline{A}$  is as above then  $\ell(\underline{A}) = k$ ). A stroll is reduced if  $F_i$  and  $F_j$  are never contained in the same hyperplane for  $i \ne j$ , i.e. if our stroll "never crosses the same reflecting hyperplane twice".

Example 2.2. Two strolls ending in the same element; one is reduced, one is not:



Remark 2.2. Starting in §3.3.5, we will redefine a stroll so that it also allows  $A_i = A_{i-1}$ . That is, a stroll is like a walk from alcove to alcove, where one might pause to admire the scenery. For the rest of this chapter, however,  $A_i \neq A_{i-1}$ .

An expression  $\underline{x} = (s_1, s_2, \dots, s_m)$  determines a stroll  $\underline{A}(\underline{x})$  via

$$\underline{A}(\underline{x}) := (A_0 = \Delta, A_1 = s_1 \Delta, A_2 = s_1 s_2 \Delta, \dots, A_k = s_1 s_2 \dots s_k \Delta).$$

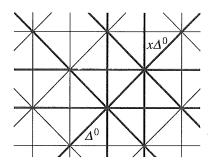
(Obviously  $\Delta$  and  $s\Delta$  meet in a codimension 1 face, and hence so do  $s\Delta$  and  $s\Delta$  for any  $s\Delta$  for any  $s\Delta$  The following proposition tells us that (reduced) expressions and (reduced) strolls are essentially the same thing:

{2\_prop:length}

**Proposition 2.1.** An expression  $\underline{x}$  for  $x \in W$  is reduced if and only if the corresponding stroll  $\underline{A}(\underline{x})$  is reduced. Moreover, we have

$$\ell(x) = \#\{\alpha \in \Phi \mid \alpha \text{ separates } \Delta^0 \text{ and } x\Delta^0\}.$$

*Example 2.3.* The geometric meaning of  $\ell(x)$ :

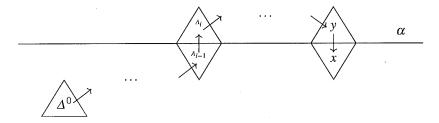


Proof. Let us temporarily define

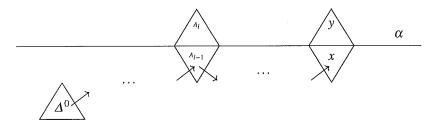
$$\ell'(x) := \#\{\alpha \in \Phi \mid \alpha \text{ separates } \Delta^0 \text{ and } x\Delta^0\}.$$

We will argue by induction on  $\ell(x)$  that  $\ell(x) = \ell'(x)$  and that any reduced expression for x yields a reduced stroll. Let  $\underline{x} = (s_1, \dots, s_k)$  denote a reduced expression for x and let  $\underline{y} = (s_1, \dots, s_{k-1})$ . Then  $\underline{y}$  is a reduced expression for  $y = s_1 \dots s_{k-1}$  (an expression of length < k - 1 for  $\underline{y}$  would yield an expression of length < k for x,

contradicting  $\ell(x) = k$ ). Thus we can apply induction to conclude that  $\ell(y) = \ell'(y)$  and that  $\underline{A}(\underline{y})$  crosses k-1 distinct hyperplanes. Now consider  $\underline{A}(\underline{x})$ . Either  $\ell(x) = \ell'(x)$  or the hyperplane  $\alpha$  crossed from  $y\Delta^0$  to  $x\Delta^0$  has already been crossed in  $\underline{A}(y)$ :



Let  $A_{i-1}$  and  $A_i$  with i < k be two alcoves where this hyperplane is crossed earlier. Then  $(s_1, \ldots, s_{i-1}, s_{i+1}, \ldots, s_{k-1})$  is an expression for x which is shorter than k. The corresponding stroll is obtained by reflecting the stroll between i and k-1 in the hyperplane  $\alpha$ :



This contradicts the fact that  $\ell(x) = k$ . Hence  $\ell(x) = \ell'(x)$  and we are done.

**Corollary 2.1.**  $x\Delta = \Delta$  if and only if x = id.

*Proof.* If  $x\Delta = \Delta$  then x is of length zero in the generators, and hence x = id.

Combining this result with Lemma 2.1 yields:

**Corollary 2.2.**  $\Delta$  is a fundamental domain for the W-action on V.

In particular the map  $x \mapsto x\Delta$  is a bijection. We can use this bijection to identify W and  $\overline{\mathscr{A}}$ . This is particularly useful as it allows us to deduce properties of W via the geometry of V and its decomposition into the sets  $\overline{\mathscr{A}}$ .

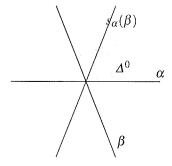
Exercise 2.3. Modify the proof of Proposition 2.1 to prove the Exchange Condition and the Deletion Condition for W (see  $\S1.2.3$ ).

## 2.4 The Coxeter presentation

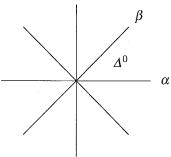
Suppose that  $\alpha$  and  $\beta$  belong to  $\Phi_{\Delta}$  (i.e.  $\alpha$  and  $\beta$  constitute walls of  $\Delta$ ).

**Lemma 2.3.** If  $\alpha$  and  $\beta$  intersect, then they do so at an angle  $\leq \pi/2$ . Moreover, this angle is of the form  $\pi/m$  for some  $m \in \mathbb{Z}_{\geq 0}$ .

*Proof.* Suppose for contradiction that  $\alpha$  and  $\beta$  intersect at an angle  $> \pi/2$ . Then reflecting  $\beta$  in the hyperplane  $\alpha$  would yield a hyperplane in the interior of  $\Delta$ , which is a contradiction:



To see the second claim is a piece of cake (by properness, the cake is cut into finitely many pieces):



If s and t denote the reflections in the hyperplanes  $\alpha, \beta \in \Phi_{\Delta}$  then we define

$$m_{st} := egin{cases} m ext{ (of the previous lemma)} & ext{if } lpha ext{ and } eta ext{ meet,} \ & ext{if } lpha ext{ and } eta ext{ do not meet.} \end{cases}$$

The composition of two reflections in distinct, parallel hyperplanes is a non-trivial translation. Meanwhile, the composition of two reflections in hyperplanes meeting at an angle of  $\pi/m$  is a rotation through  $2\pi/m$ . Hence:

**Lemma 2.4.** For s,t as above, the order of  $st \in W$  is  $m_{st}$ .

We have established the easy part (i.e. that the relations are satisfied) of the following fundamental theorem:

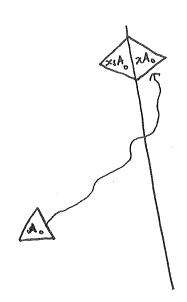
**Theorem 2.1.** W admits the following "Coxeter" presentation:

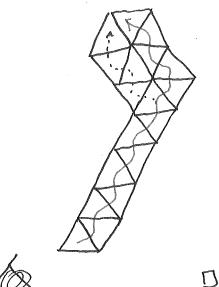
$$W = \langle s \in S \mid s^2 = \text{id } for \ all \ s \in S, (st)^{m_{st}} = \text{id } for \ all \ distinct } s, t \in S \rangle.$$

Then here exists & 1 \( \) \(

(4)

Proof:





Delivition of le extends to West:

l: Wext -> 730

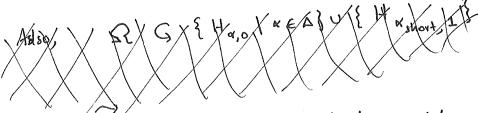
x = # {HE Je | H separates of and x A o }.

 $\Omega := \text{"length zero elements"} = \ell'(0) = \{x \mid x d_0 = d_0\}.$ 

Vertices et a Ao are {0, 81, 92,..., Brank }

and  $\Omega$   $\subset$  Sym ( {0,0,,..., oranh}) hence it is

a limite group.



SIG walls of lundamental alcore

→ St G alhie Dymlin diagram.

Jemma:

Proof: (1) Row lettous loom simple transitivity

- 2) W normal in West explained above.
- 3 Wext = ( Q) W): take x & Wext. Because W is Wansitire on alcores, 2 yEW s.t. oxy preserves do.

Hence 25' & R.

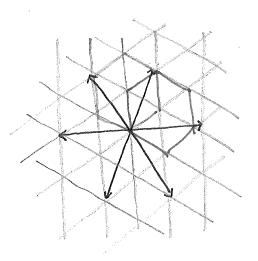
 $\Box$ 

Some examples of Q:



& s.t & presence A.

(2) A2



Sasoft & preserves sto.



B2: S= 2/22,

G( = Q= {1)

$$B = \langle T_{x}, \not \mid x \in W \mid T_{x}T_{y} = T_{xy} \text{ if } \ell(x + y) = \ell(x) + \ell(y) \rangle$$

$$B_{e} = \langle T_{x}, x \in W_{ext} \mid \dots \rangle.$$

Seel a presentation for B, Be ahin to the lattice presentation of W, We.

For 
$$\lambda \in \mathcal{X}$$
 write  $\lambda = \mathcal{X} - \mathcal{M}$ ,  $\mathcal{X}, \mathcal{M} \in \mathcal{X}_{+}$  (dominant).

Proof: # of hyperplanes  $H_{or,m}$  for  $m \in \mathbb{Z}$  separating Eg and A + Eg.

For 
$$\Delta \in \mathfrak{I}_{+}$$
,

 $\alpha \in \mathfrak{I}_{+}$ 
 $\langle \Delta, \omega^{*} \rangle$ .

Hence: 
$$\ell(\ell_{\Delta}) = \sum_{\alpha \in \Phi_{+}} \langle \gamma, \omega \gamma \rangle = 2 \langle \gamma, g^{\alpha} \rangle.$$

In pathicular, if  $\lambda$ ,  $\chi \in \mathcal{X}_{+}$ ,  $\xi = T_{\xi \lambda} T_{\xi \mu} = T_{\xi \lambda + \mu \nu}$  because

$$\ell(t_{\lambda}t_{\mu}) = \ell(t_{\lambda}t_{\mu})$$

$$= 2(\lambda t_{\mu}, s^{\nu})$$

$$= 6 \cdot \ell(t_{\lambda}) + \ell(t_{\mu}).$$

Hence: 1 C> Be via 2 - Tex.

For 
$$\Lambda \in \mathcal{K}$$
 write  $\Lambda = \Lambda' - \rho \Lambda''$ ,  $\Lambda', \Lambda'' \in \mathcal{K}_+$ .

Ts seSp

On SEX

TsTe ... = TeTs ...

mst mst

(huite braid relations)

On On = On+m

(lattice part)

ON OTS = TOOR

if  $\langle \Lambda, \alpha'_{s} \rangle = 0$ .

 $\mathcal{O}_{\lambda} T_{s}^{-1} = T_{s} \mathcal{O}_{\lambda-\alpha} \quad \text{if} \quad \langle \lambda, \alpha, \rangle = 1.$